

Minimal Exploration in Structured Stochastic Bandits

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1. Our Contribution

We investigate the stochastic Multi-Armed Bandit with finitely many arms and generic structure. We provide a generic regret lower bound and design OSSB, a generic algorithm that is asymptotically optimal for any structured MAB. We further provide a finite time analysis of OSSB.

2. Model

We consider the most general model for a stochastic structured MAB.

- The set of arms \mathcal{X} is finite
- Problem is parameterized by an unknown parameter $\theta \in \Theta$
- When arm $x \in \mathcal{X}$ is selected, one observes $Y(n, x) \sim \nu(\theta(x))$ with expectation $\theta(x)$
- Successive observations from arm x , $(Y(n, x))_n$ are i.i.d.
- When arm $x \in \mathcal{X}$ is selected, one receives a (not observed) reward $\mu(x, \theta)$
- The mapping $(x, \theta) \mapsto \mu(x, \theta)$ is known, θ is unknown
- The goal is design a policy π minimizing the regret when T is large, with $x^\pi(t)$ is the arm selected at round t by policy π :

$$R^\pi(T, \theta) = T \max_{x \in \mathcal{X}} \mu(x, \theta) - \sum_{t=1}^T \mathbb{E}(\mu(x^\pi(t), \theta)).$$

3. Structures Covered by Our Model

This model covers many popular bandit models of the literature.

- **Classical Bandits:** the parameter set is $\Theta = [0, 1]^{|\mathcal{X}|}$, the observation distribution $\nu(\theta(x))$ is any bounded distribution with mean $\theta(x)$, and for all $x \in \mathcal{X}$ the reward equals the mean observation $\mu(x, \theta) = \theta(x)$.
- **Linear Bandits:** the set of arms \mathcal{X} is a finite subset of \mathbb{R}^d ; the parameter set Θ is the set of θ such that $\theta(x) = \langle \phi, x \rangle$ for all $x \in \mathcal{X}$ for some $\phi \in \mathbb{R}^d$; the observation distribution $\nu(\theta(x))$ is a Gaussian distribution with unit variance and mean $\theta(x)$, and the reward equals the mean observation $\mu(x, \theta) = \theta(x)$.
- **Dueling Bandits:** The set of arms \mathcal{X} is $\{1, \dots, d\}^2$, and arms are $x = (i, j)$. Parameter θ is a preference matrix such that $\theta(i, j) = 1 - \theta(j, i)$, and $\theta(i, i) = \frac{1}{2}$, and there exists i^* (Condorcet winner) such that $\min_{i \neq i^*} \theta(i^*, i) > \frac{1}{2}$, the observation distribution $\nu(a)$ is the Bernoulli distribution with mean a ; the rewards are $\mu((i, j), \theta) = \frac{1}{2}(\theta(i^*, i) + \theta(i^*, j) - 1)$. Note: the best arm is (i^*, i^*) and has zero reward.
- **Lipschitz Bandits:** The set of arms \mathcal{X} is a finite metric space endowed with a distance ℓ . The mapping $x \mapsto \theta(x)$ is Lipschitz continuous with respect to ℓ , and the set of parameters is:

$$\Theta = \{\theta : |\theta(x) - \theta(y)| \leq \ell(x, y) \quad \forall x, y \in \mathcal{X}\}.$$

the reward equals the mean observation $\mu(x, \theta) = \theta(x)$.

- **Unimodal Bandits:** The set of arms \mathcal{X} is $\mathcal{X} = \{1, \dots, |\mathcal{X}|\}$, and finite and the set of parameters Θ is the set of unimodal function i.e. $x \mapsto \theta(x)$ is unimodal: it is strictly increasing on $\{1, \dots, x^*\}$ and strictly decreasing on $\{x^*, \dots, |\mathcal{X}|\}$. The reward equals the mean observation $\mu(x, \theta) = \theta(x)$.

4. Regret Lower Bounds

Assumption 1. The optimal arm $x^*(\theta)$ is unique.

Theorem 1. Let $\pi \in \Pi$ be a uniformly good algorithm. For any $\theta \in \Theta$, we have:

$$\liminf_{T \rightarrow \infty} \frac{R^\pi(T, \theta)}{\ln T} \geq C(\theta), \quad (1)$$

where $C(\theta)$ is the value of the optimization problem:

$$\underset{\eta(x) \geq 0, x \in \mathcal{X}}{\text{minimize}} \sum_{x \in \mathcal{X}} \eta(x) (\mu^*(\theta) - \mu(x, \theta)) \quad (2)$$

$$\text{s.t.} \sum_{x \in \mathcal{X}} \eta(x) D(\theta, \lambda, x) \geq 1, \quad \forall \lambda \in \Lambda(\theta), \quad (3)$$

where $D(\theta, \lambda, x)$ is the Kullback-Leibler divergence between $\nu(\theta(x))$ and $\nu(\lambda(x))$ and:

$$\Lambda(\theta) = \{\lambda \in \Theta : D(\theta, \lambda, x^*(\theta)) = 0, x^*(\theta) \neq x^*(\lambda)\}. \quad (4)$$

is the set of parameters λ where the optimal arm $x^*(\lambda)$ is different from $x^*(\theta)$ and cannot be distinguished from θ by sampling $x^*(\theta)$.

5. Bounds for Specific Structures

The regret lower bound covers previously known lower bounds for specific structured bandits. Also, the solution of (2)-(3) is often tractable.

- **Classical bandits:** (Lai, 1985) the solution of (2)-(3) is:

$$c(x, \theta) = \frac{1}{d(\theta(x), \theta(x^*))}.$$

- **Linear bandits:** for Gaussian rewards (Lattimore et al. 2016), (2)-(3) is equivalent to the convex program:

$$\underset{\eta(x) \geq 0, x \in \mathcal{X}}{\text{minimize}} \sum_{x \in \mathcal{X}} \eta(x) (\theta(x^*) - \theta(x))$$

$$\text{s.t.} x^\top \text{inv} \left(\sum_{z \in \mathcal{X}} \eta(z) z z^\top \right) x \leq \frac{(\theta(x^*) - \theta(x))^2}{2}, \quad \forall x \neq x^*.$$

- **Lipschitz bandits:** for Bernoulli rewards (Magureanu et al. 2014), (2)-(3) is equivalent to a linear program ($|\mathcal{X}|$ variables and $2|\mathcal{X}|$ constraints):

$$\underset{\eta(x) \geq 0, x \in \mathcal{X}}{\text{minimize}} \sum_{x \in \mathcal{X}} \eta(x) (\theta(x^*) - \theta(x))$$

$$\text{s.t.} \sum_{z \in \mathcal{X}} \eta(z) d(\theta(z), \max\{\theta(z), \theta(x^*) - \ell(x, z)\}) \geq 1, \quad \forall x \neq x^*.$$

- **Dueling bandits:** if there exists a Condorcet winner i^* (Komyama, 2016), the solution of (2)-(3) is: (where $x = (i, j)$)

$$c(x, \theta) = \mathbb{1} \left\{ j \in \arg \min_{j'} \frac{\mu((i, j'), \theta)}{d(\theta(i, j'), 1/2)} \right\} \frac{1}{d(\theta(i, j), 1/2)}.$$

- **Unimodal bandits:** (Combes et al. 2014) the solution of (2)-(3) is:

$$c(x, \theta) = \frac{\mathbb{1}\{|x - x^*| = 1\}}{d(\theta(x), \theta(x^*))}.$$

6. The OSSB Algorithm

OSSB(ϵ, γ) Pseudocode.

$s(0) \leftarrow 0, N(x, 1), m(x, 1) \leftarrow 0, \forall x \in \mathcal{X}$ {Initialization}
for $t = 1, \dots, T$ do

Compute the optimization problem (2)-(3) solution $(c(x, m(t)))_{x \in \mathcal{X}}$
 where $m(t) = (m(x, t))_{x \in \mathcal{X}}$

if $N(x, t) \geq c(x, m(t))(1 + \gamma) \ln t, \forall x$ **then**

$s(t) \leftarrow s(t - 1)$

$x(t) \leftarrow x^*(m(t))$ {Exploitation}

else

$s(t) \leftarrow s(t - 1) + 1$

$\bar{X}(t) \leftarrow \arg \min_{x \in \mathcal{X}} \frac{N(x, t)}{c(x, m(t))}$

$\underline{X}(t) \leftarrow \arg \min_{x \in \mathcal{X}} N(x, t)$

if $N(\underline{X}(t), t) \leq \epsilon s(t)$ **then**

$x(t) \leftarrow \underline{X}(t)$ {Estimation}

else

$x(t) \leftarrow \bar{X}(t)$ {Exploration}

end if

end if

 Play $x(t)$ and update statistics.

end for

OSSB(ϵ, γ) is provably asymptotically optimal (we give a finite time analysis).

Assumption 2. For all x , the mapping $(\theta, \lambda) \mapsto D(x, \theta, \lambda)$ is continuous at all points where it is not infinite.

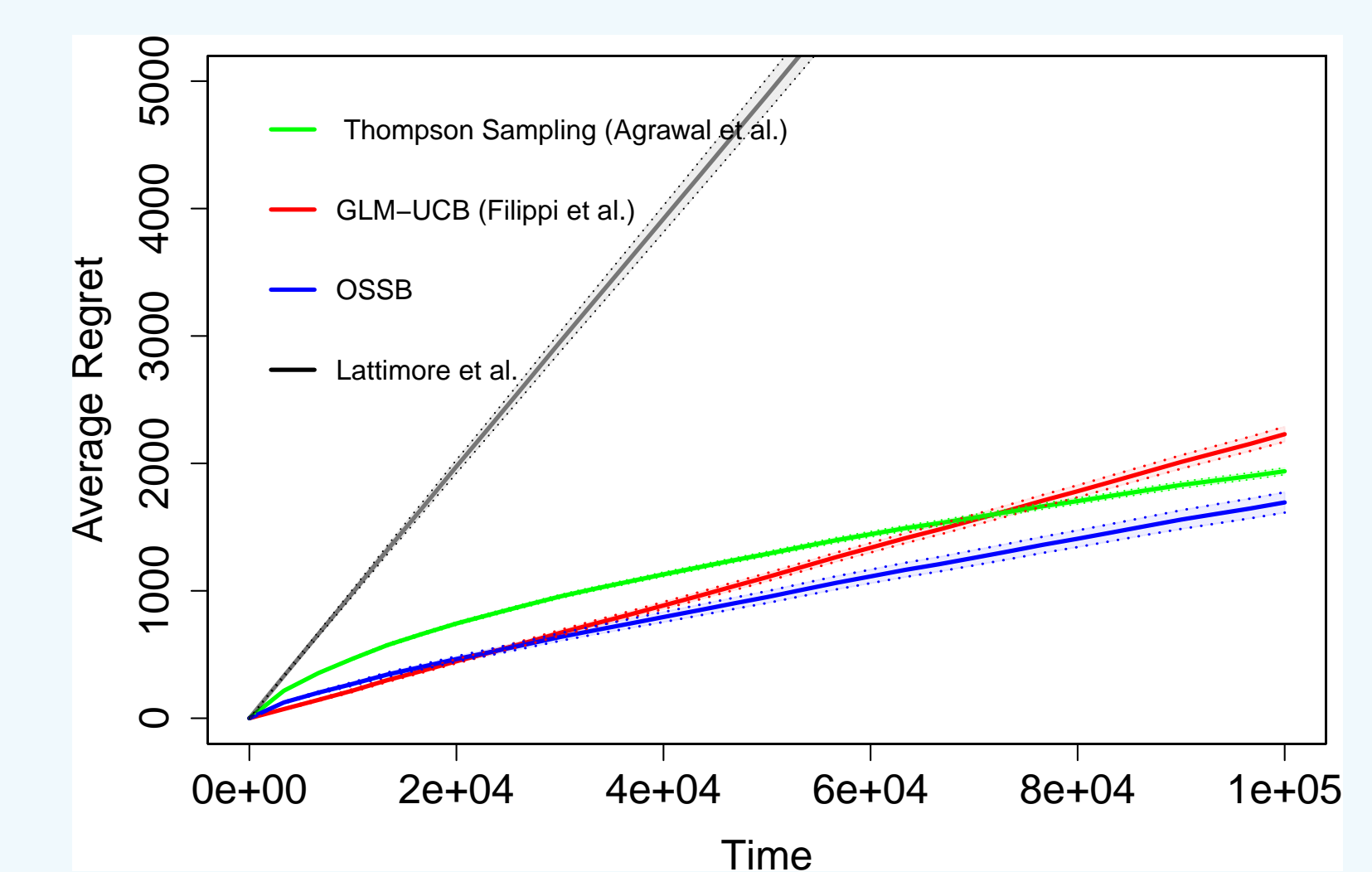
Assumption 3. For all x , the mapping $\theta \rightarrow \mu(x, \theta)$ is continuous.

Theorem 2. Under Assumptions 1, 2 and 3, for Bernoulli and Gaussian observation, then under the algorithm $\pi = \text{OSSB}(\epsilon, \gamma)$ with $\epsilon < \frac{1}{|\mathcal{X}|}$ we have:

$$\limsup_{T \rightarrow \infty} \frac{R^\pi(T)}{\ln T} \leq C(\theta) F(\epsilon, \gamma, \theta),$$

with F a function such that for all θ , we have $F(\epsilon, \gamma, \theta) \rightarrow 1$ as $\epsilon, \gamma \rightarrow 0$.

7. Numerical Results



Performance of OSSB(0,0) vs. state-of-the-art algorithms.

- We consider linear bandits with 81 arms and 10 random instances.
- OSSB works well in finite time (competitive with the state of the art).
- Since OSSB is more complex to implement than other algorithms, reducing its complexity is an interesting topic of future research.